

**MATH 521A: Abstract Algebra**  
Homework 1 Solutions

1. Prove that  $(-\mathbb{N}_0)$ , the set of nonpositive integers, is well-ordered.

Method 1: We know that  $\mathbb{Z}$  is well-ordered by  $\leq$ , by a theorem from class. Since  $(-\mathbb{N}_0) \subseteq \mathbb{Z}$ , we apply the lemma from class and conclude that  $(-\mathbb{N}_0)$  is also well-ordered by  $\leq$ .

Method 2: We define a “backwards” order via:  $a < b$  if  $|a| < |b|$  (i.e.  $a > b$ ). Let  $S \subseteq (-\mathbb{N}_0)$ . The image of  $S$  under the absolute value map is in  $\mathbb{N}_0$ , which is well-ordered. Hence there is some element  $t$  in that image that is minimal. But then  $|t| < |x|$  for all  $x \in S$ , i.e.  $t < x$ . So  $t$  is minimal and  $(-\mathbb{N}_0)$  is well-ordered by  $<$ .

For a set  $T$ , we say it is *inductively ordered* if there is some special  $t \in T$  and some function  $f : T \rightarrow T$  such that:

- (1) The elements  $t, f(t), f(f(t)), \dots$  are all distinct; and
- (2)  $T = \{t, f(t), f(f(t)), \dots\}$ .

2. Prove that  $\mathbb{N}_0$  is inductively ordered.

We take  $t = 0$  and  $S(x) = x + 1$ . Then  $\{t, f(t), f(f(t)), \dots\} = \{0, 1, 2, \dots\} = \mathbb{N}_0$ .

3. Prove that if a set is inductively ordered then it is well-ordered.

Each element of  $T$  is  $f^{(n)}(t)$ , for some  $n \in \mathbb{N}_0$ . We define an order on  $T$  by comparing “exponents”, i.e. via  $f^{(n)}(t) < f^{(m)}(t)$  if  $n < m$ . For  $S \subseteq T$ , the exponents of the elements of  $S$  are a subset of  $\mathbb{N}_0$ , and hence have a minimal element  $n^*$ . Now for any  $s \in S$ , either  $s = f^{(n^*)}(t)$  or  $f^{(n^*)}(t) < s$ , so  $f^{(n^*)}(t)$  is minimal in  $S$ . Thus  $T$  is well-ordered by  $<$ .

4. Prove that the square of any integer  $a$  is either of the form  $4k$  or of the form  $4k + 1$  for some integer  $k$ .

Let  $a \in \mathbb{Z}$ . By the division algorithm, there are integers  $q, r$  such that  $a = 4q + r$ , with  $0 \leq r < 4$ . Squaring, we get  $a^2 = (4q + r)^2 = 16q^2 + 8qr + r^2 = 4(4q^2 + 2qr) + r^2$ . If  $r = 0$  or  $r = 1$ , we are done. If  $r = 2$  then  $a^2 = 4(4q^2 + 2qr + 1) + 0$ , and if  $r = 3$  then  $a^2 = 4(4q^2 + 2qr) + 9 = 4(4q^2 + 2qr + 2) + 1$ .

5. Prove the *Backwards Division Algorithm*: Let  $a, b$  be integers with  $b > 0$ . Then there exist integers  $q, r$  such that  $a = bq + r$  with  $-b < r \leq 0$ .

We mimic the proof of Thm 1.1. Set  $S = \{a - bx : x \in \mathbb{Z}, a - bx \geq 1\}$ . In Thm 1.1 it was proved that there is some  $x \in \mathbb{Z}$  such that  $a - bx \geq 0$ . Adding  $b$  to both sides  $a - bx + b \geq b \geq 1$ , so  $S$  is nonempty since it contains  $a - b(x - 1)$ . Apply Well-Ordering Axiom to get minimal  $r' \in S$ . We have, for some integer  $q'$ ,  $r' = a - bq' \geq 1$ . We subtract  $b$  from both sides to get  $r' - b = a - b(q' + 1) \geq 1 - b$ . Set  $r = r' - b$  and  $q = q' + 1$ . We have  $r = a - bq$ , or  $a = bq + r$ . If  $r > 0$  then  $r' > b$  so in fact  $r'$  was not minimal in  $S$  since  $r' - b \in S$ . Hence  $-b < r \leq 0$ . We were not asked to prove uniqueness, but it can be done similarly to the proof of Thm 1.1.

6. Let  $a, b \in \mathbb{N}$  with  $a|b$ . Prove that  $a \leq b$ .

Method 1: We use  $x \geq y$  if  $x - y \geq 0$ . There is some  $c \in \mathbb{N}$  with  $b = ac$ . Hence  $b - a = ac - a = a(c - 1)$ . Since  $a, c - 1$  are each in  $\mathbb{N}_0$ , their product is in  $\mathbb{N}_0$  and so  $b \geq a$ .

Method 2: We use  $x \geq y$  if  $\frac{x}{y} \geq 1$ . There is some  $c \in \mathbb{N}$  with  $b = ac$ . Hence  $\frac{b}{a} = c \in \mathbb{N}$ . Hence  $\frac{b}{a} \geq 1$ , so  $b \geq a$ .

7. Let  $a, b$  be nonzero integers with  $a|b$  and  $b|a$ . Prove that  $a = \pm b$ .

Since  $a|b$  there is some integer  $c$  with  $b = ca$ . Since  $b|a$  there is some integer  $f$  with  $a = fb$ . Substituting, we get  $b = c(fb)$  and so  $1 = cf$ . Suppose for the moment that  $c, f$  are both positive. We have  $c|1$ , so by exercise 6,  $c \leq 1$ . Thus  $c = 1 = f$ . If  $c, f$  are not both positive, they are both negative. However  $|c| \cdot |f| = 1$  so again  $|c|$  is a natural number dividing 1, and thus  $|c| = 1$ , so  $c = -1 = f$ . Hence  $c = \pm 1$  and so  $b = \pm a$ .

8. Let  $a, b \in \mathbb{Z}$ , not both zero, and let  $d = \gcd(a, b)$ . Prove that  $d$  divides each element of  $S = \{am + bn : m, n \in \mathbb{Z}\}$ .

We have  $a = da', b = db'$  for some integers  $a', b'$ . We have  $am + bn = da'm + db'n = d(a'm + b'n)$ , so  $d|(am + bn)$ .

9. Use the Euclidean Algorithm to find  $\gcd(175, 630)$  and to express this as a linear combination of 175, 630.

Step 1:  $630 = 3 \cdot 175 + 105$ . Step 2:  $175 = 1 \cdot 105 + 70$ . Step 3:  $105 = 1 \cdot 70 + 35$   
Since  $35|70$  we know 35 is the gcd. Now we reverse, back-substituting as we go.  
 $35 = 1 \cdot 105 - 1 \cdot 70 = 1 \cdot 105 - 1 \cdot (1 \cdot 175 - 1 \cdot 105) = 2 \cdot 105 - 1 \cdot 175 = 2 \cdot (630 - 3 \cdot 175) - 1 \cdot 175 = 2 \cdot 630 - 7 \cdot 175$ .

10. Prove that  $\gcd(a, b) = \gcd(a, b + at)$ , for every  $t \in \mathbb{Z}$ .

Method 1: Fix  $t$  and for convenience, set  $d = \gcd(a, b), c = \gcd(a, b + at)$ . Since  $d|a$  and  $d|b$ , there are integers  $a', b'$  with  $a = da', b = db'$ . So  $b + at = db' + da't = d(b' + a't)$ , so  $d|(b + at)$  and hence  $d|c$  by Cor. 1.3. Similarly, there are  $a'', f$  such that  $a = a''c, b + at = fc$ . Hence  $(b + at) - at = fc - a''ct = c(f - a''t)$ . Hence  $c|b$  and so  $c|d$ . By exercise 7,  $c = \pm d$ , so  $c = d$  since gcd's are always positive.

Method 2: We have  $\gcd(a, b)$  dividing  $a$ , and also (by Exercise 8)  $\gcd(a, b + at)$ , since  $b + at$  is a linear combination of  $a, b$ . But we can do this the other way, since  $a = (-t)a + 1(b + at)$ , so  $\gcd(a, b + at)$  divides both  $a$  and  $b$ , so  $\gcd(a, b + at)$  divides  $\gcd(a, b)$ . Now apply exercise 7 as in method 1.